

1. What is the value of y after both of the following operations?

```
x = x ^ (~y);  
y = y ^ x;
```

**~x**

$y = y \wedge (x \wedge (\sim y)) \rightarrow (y \wedge \sim y) \wedge x \rightarrow 1s \wedge x == \sim x$

After you plug in x, you can use the commutative and associative properties of XOR and do  $y \wedge \sim y$  first which results in all 1s. x XORed with 1s flips its bits, thus  $\sim x$

Say  $x = 0111$  and  $y$  is  $1010$   
 $0111 \wedge 0101 = 0010$   
 $1010 \wedge 0010 = 1000$  which is  $\sim x$

2. Given the following declarations, do the statements below always evaluate to true?

```
int x = foo();  
int y = bar();  
unsigned ux = cookie();
```

a.

$x > ux \implies (\sim x + 1) < 0$

**FALSE**

Consider  $x = -1$ .

- The binary is all 1s, thus when you do  $\sim$ (all 1s) it becomes all 0s.
  - Adding the 1 makes the value positive.

This is true for all negative x values since the sign bit will always be flipped to 0.

- So the 'it follows' is not true for all  $x > ux$ .

b.

$ux - 2 \geq -2 \implies ux \leq 1$

**TRUE**

If  $ux$  is 0

- it is comparing the unsigned values of -2 and -2, which are equal.

If  $ux$  is 1

- it is comparing the unsigned values of -1 and -2, which are  $U_{max}$  vs  $U_{max} - 1$ .  
2,3, etc
  - aren't true and  $ux$  can not be a negative value.
- So, it follows that  $ux$  must be 0 or 1.

c.

$$(x^y)^x == (x+y)^{(x+y)^y}$$

TRUE

Notice that both sides are of the form  $(A^y)^A$

- For the left hand side,  $A = x$
- For the right hand side,  $A = x+y$

$(A^y)^A$  is equivalent to  $y$

- Thus, the equivalence simplifies to  $y == y$
- Both sides of the equivalence are equal

d.

$$(x < 0) \ \&\& \ (y < 0) == (x + y) < 0$$

FALSE

Say  $x == INTMin$  and  $y == INTMin$ .

- $(x+y)$  would overflow.

- |                                     |           |               |
|-------------------------------------|-----------|---------------|
| 3. <code>char** apple[5][9];</code> | 360 bytes | $(8 * 5 * 9)$ |
| <code>char* banana[1][9];</code>    | 72 bytes  | $(8 * 1 * 9)$ |
| <code>char strawberry[4][2];</code> | 8 bytes   | $(1 * 4 * 2)$ |

How many bytes of space would these declarations require?

4. Consider the following struct:

```
typedef struct {
    char first;
    int second;
    short third;
} stuff;
```

Say we are debugging an application in execution using gdb on a 64-bit, little-endian architecture. The application has a variable called array - defined as:  
`stuff array[2][2];`

Using gdb we find the following information at a particular stage in the application:

```
[(gdb) p &array  
$1 = (stuff (*)[2][2]) 0x7fffffff020
```

And:

```
[(gdb) x/48xb 0x7fffffff020  
0x7fffffff020: 0x61 0x00 0x00 0x00 0x08 0x00 0x00 0x00  
0x7fffffff028: 0x02 0x00 0x00 0x00 0x62 0x00 0x00 0x00  
0x7fffffff030: 0x64 0x00 0x00 0x00 0x04 0x00 0x00 0x00  
0x7fffffff038: 0x63 0x04 0x40 0x00 0xed 0x03 0x00 0x00  
0x7fffffff040: 0xc8 0x00 0xff 0xff 0x64 0x7f 0x00 0x00  
0x7fffffff048: 0x17 0xa6 0x00 0x00 0xe1 0x00 0x00 0x00
```

What is the value of  
`array[1][0].second`

At this particular stage of the application?  
i.e. what would be returned from the statement:

```
printf("%d\n", array[1][0].second);
```

### 1005

Because of alignment, each object of type “stuff” is 12 bytes.

Due to how arrays are stored in memory,

- The array is stored as:  
`array[0][0], array[0][1], array[1][0], array[1][1]`

From the gdb output, we can tell that the array starts at `0x7fffffff020`

- `array[1][0]` is `0x7fffffff038` to `0x7fffffff043`
  - Note: this is in hex, so  $0x7fffffff038 + 8 = 0x7fffffff040$

`second` is an integer, and is the 5th to 8th byte of an object of type “stuff”

- These are bytes `0x7fffffff03c` to `0x7fffffff03f`
- They have the values `0xed`, `0x03`, `0x00`, `0x00`
- Since this system is little endian, the value is `0x000003ed`
  - This is equivalent to 1005

5. The following is part of the result of the command `'objdump -d'` on an executable

```

0000000004006dd <IronMan>:
 4006dd:    55                push   %rbp
 4006de:    48 89 e5         mov    %rsp,%rbp
 4006e1:    89 7d ec         mov    %edi,-0x14(%rbp)
 4006e4:    8b 45 ec         mov    -0x14(%rbp),%eax
 4006e7:    c1 e0 04        shl   $0x4,%eax
 4006ea:    89 45 fc         mov    %eax,-0x4(%rbp)
 4006ed:    8b 45 fc         mov    -0x4(%rbp),%eax
 4006f0:    5d                pop    %rbp
 4006f1:    c3                retq

```

Say the declaration for the function IronMan was:

```
int IronMan(int scraps);
```

Given that the integer 23 was passed into the function, what is the return value?

368

After instructions 0x4006e1 and 0x4006e4, the input (which was stored in %rdi) is now stored in %eax

Instructions 0x4006e7 then shifts %eax to the left by 4

- This is equivalent to multiply by  $2^4$ , which is 16  
 $23 * 16 = 368$

6. The following is a continuation from the previous problem:

```
0000000000400721 <Hulk>:
400721: 55                push   %rbp
400722: 48 89 e5          mov    %rsp,%rbp
400725: 48 83 ec 20       sub    $0x20,%rsp
400729: 48 89 7d e8       mov    %rdi,-0x18(%rbp)
40072d: 48 8b 45 e8       mov    -0x18(%rbp),%rax
400731: 48 89 c7          mov    %rax,%rdi
400734: e8 27 fe ff ff   callq 400560 <atoi@plt>
400739: 89 45 fc          mov    %eax,-0x4(%rbp)
40073c: 8b 45 fc          mov    -0x4(%rbp),%eax
40073f: 89 c7            mov    %eax,%edi
400741: e8 97 ff ff ff   callq 4006dd <IronMan>
400746: 89 45 f8          mov    %eax,-0x8(%rbp)
400749: 81 7d f8 8f 01 00 00  cmpl  $0x18f,-0x8(%rbp)
400750: 7e 10            jle   400762 <Hulk+0x41>
400752: 81 7d f8 f4 01 00 00  cmpl  $0x1f4,-0x8(%rbp)
400759: 7f 07            jg   400762 <Hulk+0x41>
40075b: b8 01 00 00 00    mov    $0x1,%eax
400760: eb 05            jmp   400767 <Hulk+0x46>
400762: b8 00 00 00 00    mov    $0x0,%eax
400767: c9              leaveq
400768: c3              retq
```

Given that the function returns 1, what do we know about the value of %edi right before instruction 0x400741 is executed?

%edi is between 25 and 31

Since the function returns 1, we know that the jump instructions at 0x400750 and 0x400759 did not jump.

- From instructions 0x400749 and 0x400750
  - we know that we would have jumped if  $-0x8(\%rbp)$  was less than or equal to 0x18f
  - Thus we know  $-0x8(\%rbp)$  is greater than 0x18f, or 399
- From instructions 0x400752 and 0x400759
  - We know that we would have jumped if  $-0x8(\%rbp)$  was greater than 0x1f4
  - Thus we know  $-0x8(\%rbp)$  is less than or equal to 0x1f4, or 500
- Thus we know that  $-0x8(\%rbp)$  is between 400 and 500, inclusive
  - Thus %eax is between 400 and 500, inclusive

From the previous question, we know that IronMan multiplies inputs by 16

- We also know that the function returns a value between 400 and 500 with input %rdi
- Reversing the function, we know the input must have been between 400/16 and 500/16

Thus we know that %rdi was between 25 and 31 right before the IronMan function call

7. Convert the 8-bit floating point number e7 (in hex) to decimal.

A. Convert:  $e7_{16} = 11100111_2$ .

B. Separate: 

1	110	0111
---	-----	------

C. Mantissa: 1.0111

D. Exponent:  $110_2 = 6_{10}; 6 - 3 = 3$ .

E. De-normalize:  $1.0111_2 \times 2^3 = 1011.1$

F. Convert:

<i>Exponents</i>	$2^3$	$2^2$	$2^1$	$2^0$	$2^{-1}$	
<i>Place Values</i>	8	4	2	1	0.5	
<i>Bits</i>	1	0	1	1	1	
<i>Value</i>	8			+ 2	+ 1	+ 0.5 = 11.5

G. Sign: negative.

Result: e7 is -11.5

b) Convert the number -15.5 into floating point, using 1 sign bit, 3 exponent bits, and 4 mantissa bits.

$$-15.5 = -1111.1 \text{ } (-15 + 1^{-1})$$

$$-1111.1 = -1.1111 \times 2^3 \text{ (where } -1.1111 \text{ is base 2)}$$

We have 3 exponent bits, so we calculate our bias as  $2^{(3 - 1)} - 1 = 3$

Our sign bit is 1, because our number is negative.

Our exponent bits are 110, because we have an exponent of 3, and a bias of 3.

Our mantissa bits are 1111, because that is our decimal after the .

Our final result:

1110 1111