

Worksheet 2 Solutions

1. `mov` vs `lea` - describe the difference between the following:

```
movl (%rdx), %rax
```

```
leal (%rdx), %rax
```

```
movl rdx, %rax
```

`movl` takes the **contents** of what's stored in register `%rdx` and moves it to `%rax`. `leal` computes the load effective **address** and stores it in `%rax`. `leal` analogous to returning a pointer, whereas `movl` is analogous to returning a dereferenced pointer.

2. Invalid `mov` Instructions

Explain why these instructions would not be found in an assembly program.

- | | |
|--------------------------------------|---|
| a) <code>movl %eax, %rdx</code> | destination operand has the incorrect size |
| b) <code>movb %di, 8(%rdx)</code> | instruction suffix does not match register size |
| c) <code>movq (%rsi), 8(%rbp)</code> | source and destination cannot both be memory references |
| d) <code>movw 0xFF, (%eax)</code> | <code>%eax</code> cannot be used as an address register (not 64 bits) |

3. What would be the corresponding instruction to move 64 bits of data from register `%rax` to register `%rcx`?

```
movq (%rax), %rcx
```

(important part is that you know the suffix of the MOV instruction!)

4. Operand Form Practice (see page 181 in textbook)

Assume the following values are stored in the indicated registers/memory addresses.

<u>Address</u>	<u>Value</u>	<u>Register</u>	<u>Value</u>
0x104	0x34	%rax	0x104
0x108	0xCC	%rcx	0x5
0x10C	0x19	%rdx	0x3
0x110	0x42	%rbx	0x4

Fill in the table for the indicated operands:

<u>Operand</u>	<u>Value</u>	<u>Operand</u>	<u>Value</u>
\$0x110	0x110 (immediate value)	3(%rax, %rcx)	0x19 (value in %rax is 0x104, value in %rcx is 0x5, $3 + 0x104 + 0x5 = 0x10C$, value in 0x10C is 0x19)
%rax	0x104 (value stored in %rax)	256(, %rbx, 2)	0xCC (value in %rbx is 0x4, 256 in hex is 0x100, $0x100 + (0x4 * 2) = 0x108$, value in memory address 0x108 is 0xCC)
0x110	0x42 (value stored in memory address 0x110)	(%rax, %rbx, 2)	0x19 (value in %rax is 0x104, value in %rbx is 0x4, $0x104 + (0x4 * 2) = 0x10C$, value in memory address 0x10C is 0x19)

```
(%rax)          0x34
                (%rax holds 0x104,
memory address 0x104
                holds 0x34)
```

```
8(%rax)        0x19
                (%rax holds 0x104, 8
                + 0x104 = 0x10C,
                value in memory
                address 0x10C is
                0x19)
```

```
(%rax, %rbx)   0xCC
                (value in %rax is
                0x104, value in %rbx
                is 0x4, 0x104 + 0x4
                = 0x108, value in
                memory address 0x108
                is 0xCC)
```

- \$ denotes immediates
- Note: any numbers starting with "0x" are hexadecimal numbers!!
- All of the operands can be evaluated using the specific formulas on page 181 in the textbook
- More generally, whenever you see an address of the form $D(r_b, r_i, s)$, where D is an number, r_b and r_i are registers, and s is either 1, 2, 4, or 8, you can use the following formula:

$$D + R[r_b] + R[r_i]*s$$

If D is missing, assume $D == 0$

If r_b is missing, assume $r_b == 0$

If r_i is missing, assume $r_i == 0$

If s is missing, assume $s == 1$

- For more practice, try practice problem 3.1 on page 182 of the textbook

5. Condition Codes and Jumps

Assume the addresses and registers are in the same state as in Problem 2. Does the following code result in a jump to .L2?

```
leaq (%rax, %rbx), %rdi
cmpq $0x100, %rdi
jg .L2
```

Yes.

1. First line will put $0x104+0x04 = 0x108$ into %rdi.
2. Second line sets codes according to $0x108 - 0x100$, which sets no codes.
3. Since jg is evaluated as $\sim(SF^{\wedge}OF)\&\sim ZF$ which in this case is $\sim(0^{\wedge}0)\&\sim 0 = 1\&1 = 1$.
4. So we will jump.

Students can make mistake of doing a memory read in leaq which would put $0xCC$ into %rdi. This would mean $ZF = 1$ after cmpq, which would mean jg logic evaluates to false (no jump).

Problem 6

```
int cool1(int a, int b) {
    if ( b < a )
        return b;
    else
        return a;
}

int cool2(int a, int b) {
    if ( a < b )
        return a;
    else
        return b;
}

int cool3(int a, int b) {
    unsigned ub = (unsigned) b;
    if ( ub < a )
        return a;
    else
        return ub;
}
```

Which of the functions would compile into this assembly code:

```
    movl %csi, %eax
    cmpl %eax, %edi //a-b
    jge .L4 //SF - OF
    movl %edi, %eax
.L4:  ret
```

cool2

- Arguments passed to a function is stored in the %edi, %esi, etc registers
 - %edi is a and %esi is b
- When comparing, we compare as `cmp Two One`
 - Thus the instruction `jge` is checking if %edi is greater than or equal to %eax
 - This is essentially checking if $a \geq b$, which is the else condition
- We can observe that when we do jump, %eax is not updated
 - We return b in the else case

- If we don't jump, we update %eax to %edi
 - We return a in the if case
- Thus cool2
- This question was inspired by a previous midterm